
Software Driver

OS Abstraction Middleware

Introduction

This Application Note describes the operation of the Renesas OS Abstraction middleware for Renesas microcontrollers. This document does assume that the reader has some knowledge of e² studio and CS+.

Target Device

Renesas Microcontrollers

Driver Dependencies

For OS abstraction with an embedded OS, the middleware requires the underlying OS to be within the project.

For OS abstraction without an embedded OS, the middleware requires the OSTM driver to be within the project.

List of Abbreviations and Acronyms

Abbreviation	Full Form
API	Application Programming Interface
ASCII	American Standard Code for Information Interchange
ISR	Interrupt Service Routine
OS	Operating System
OSTM	Operating System Timer Module
RTOS	Real Time Operating System

Table 1-1 List of Abbreviations and Acronyms

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1. Outline of OS Abstraction

The OS Abstraction middleware provides the user with a standardized API to operating system features for process and task control.

By using a common, consistent API for OS access, the effort involved with porting application code to different operating systems is greatly simplified. Furthermore, with OS-less OS abstraction, a common approach is provided for non-OS environments as well.

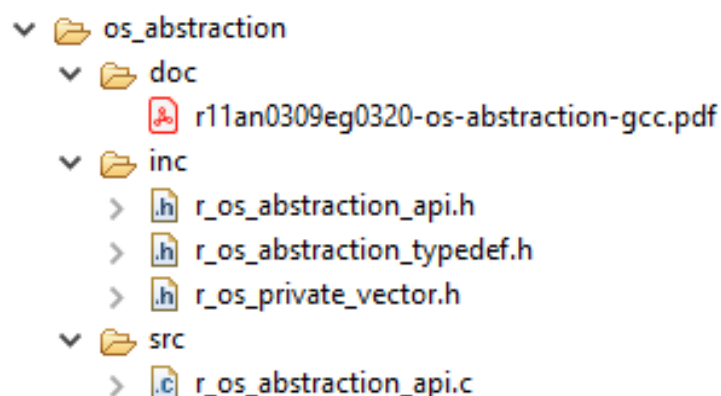
2. Description of the Middleware

The key features to configure:

- Tasks – not used in the OS Less variant of this API
- Mutexes
- Semaphores
- Memory Allocation
- Events
- Message Queues

2.1 Structure

An example of the OS abstraction file structure can be seen in the image below.



2.2 Description of each file

Each file's description can be seen in the following table.

Filename	Usage	Description
r_os_abstraction_api.h	To be included in any file which executes the OS Abstraction API	This and r_task_priority.h are the only API header files to include in application code
r_os_private_vector.h	System Configuration only	System Configuration only
r_os_abstraction_typedef.h	Included by r_os_abstraction_api.h	Defines OS abstraction data types
r_task_priority.h	Included by the application	Task priority definitions. Not required if OS-less OS abstraction is used.
r_os_abstraction_api.c	Private	The OS abstraction code implementation.

3. Example of Use

This section describes a simple example of creating a task, mutex, semaphore, event and message queue.

3.1 Create Task

```
os_task_t * p_os_task;

p_os_task = R_OS_TaskCreate("My Task", my_task_function, NULL,
    R_OS_ABSTRACTION_SMALL_STACK_SIZE, 6);

if (NULL == p_os_task)
{
    printf("Task Creation Error");
}
```

3.2 Create Mutex

```
void *p_mutex = R_OS_MutexCreate();
```

3.3 Create Semaphore

```
uint32_t my_semaphore = 0;
uint32_t count = 10u;
bool_t success;

success = R_OS_SemaphoreCreate((p_semaphore_t) &my_semaphore, count);
if (!success)
{
    printf("Semaphore Creation Error");
}
```

3.4 Create Event

```
p_event_t my_event = NULL;
bool_t success;

success = R_OS_EventCreate(&my_event);
if (!success)
{
    printf("Event Creation Error");
}
```

3.5 Create Message Queue

```
uint32_t queue_size = 10u;
bool_t success;
p_os_msg_queue_handle_t my_message_queue_handle;

success = R_OS_MessageQueueCreate(&my_message_queue_handle, queue_size);
if (!success)
{
    printf("Message Queue Creation Error");
}
```

4. Module Documentation

4.1 Detailed Description

Provides OS abstraction, use these primitives in the code base NOT direct calls to underlying OS primitives.

Provides type defines for OS abstraction.

To make efficient code re-use the identical API shall be used in both OS and OS Less applications. This file aims to abstract the Operating system (OS) awareness when creating an OS Less driver.

4.2 Known Limitations

R_OS_TaskUsesFloatingPoint should be called in the task that use FPU.

4.3 Known Implementations

RZA2M Software packages

4.4 Related modules

See also: DS_BOARD_SUPPORT, RZA1H_RSK_OSTM_DRIVER, RZA1H_RSK_LED

4.5 Macro Definition Documentation

```
#define SRC_RENESAS_APPLICATION_INC_R_OS_ABSTRACTION_API_H_
```

```
#define R_OS_ABSTRACTION_VERSION_MAJOR (1)
```

```
#define R_OS_ABSTRACTION_VERSION_MINOR (0)
```

```
#define R_OS_ABSTRACTION_UID (81)
```

```
#define R_OS_ABSTRACTION_BUILD_NUM (0)
```

Build Number of API.

Generated during customer release.

```
#define R_OS_ABSTRACTION_EV_WAIT_INFINITE (0xFFFFFFFFFUL)
```

Maximum timeout used in wait functions inside the OS abstraction module

```
#define R_OS_ABSTRACTION_INVALID_HANDLE (-1)
```

Invalid handle used in functions inside the OS abstraction module

```
#define R_OS_ABSTRACTION_TINY_STACK_SIZE (0)
```

Stack sizes, these indexes are mapped to actual sizes inside the OS abstraction module

```
#define R_OS_ABSTRACTION_SMALL_STACK_SIZE (1)
```

```
#define R_OS_ABSTRACTION_DEFAULT_STACK_SIZE (2)
```

```
#define R_OS_ABSTRACTION_LARGE_STACK_SIZE (3)
```

```
#define R_OS_ABSTRACTION_HUGE_STACK_SIZE (4)
```

```
#define R_OS_ABSTRACTION_MAX_TASK_NAME_SIZE (24)
```

```
#define R_OSFREE_MAX_MUTEXES (32)
```

Max number of simultaneous mutexes available. Adjust to suit application

```
#define R_OSFREE_MAX_EVENTS (32)
```

Max number of simultaneous events available. Adjust to suit application

```
#define R_OS_ABSTRACTION_OSTM_RESOURCE ("\\\\.\\ostm_reserved")
```

```
#define R_OS_MS_TO_SYSTICKS(n) (n)
```

```
#define R_OS_SYSTICKS_TO_MS(n) (n)
```

4.6 Function Documentation

`bool_t R_OS_AbstractionLayerInit (void)`

Function to configure critical resources for the connected OS or scheduler.

Return values:

<i>true</i>	if there were no errors when initialising the OS Abstraction Layer.
<i>false</i>	if there errors when initialising the OS Abstraction Layer.

`bool_t R_OS_AbstractionLayerShutdown (void)`

Function to release critical resources for the connected OS or scheduler.

Return values:

<i>true</i>	if there were no errors when closing the OS Abstraction Layer.
<i>false</i>	if there errors when closing the OS Abstraction Layer.

`void R_OS_AssertCalled (volatile const char * p_file, volatile uint32_t line)`

Generic error handler, enters forever loop but allows debugger to step out..

Parameters:

in	<i>file</i>	file in which the error occurred.
in	<i>line</i>	line where the error occurred.

Return values:

<i>NONE.</i>	
--------------	--

`void R_OS_EnterCritical (void)`

Enter critical area of code - prevent context switches.

OS Abstraction R_OS_EnterCritical Function

`bool_t R_OS_EventCreate (pp_event_t pp_event)`

Create an event object for inter-task communication.

Parameters:

in	<i>pp_event</i>	Pointer to an associated event.
----	-----------------	---------------------------------

Returns:

The function returns TRUE if the event object was successfully created. Otherwise, FALSE is returned

`void R_OS_EventDelete (pp_event_t pp_event)`

Delete an event, freeing any associated resources.

Parameters:

in	<i>pp_event</i>	Pointer to an associated event.
----	-----------------	---------------------------------

Returns:

none

`e_event_state_t R_OS_EventGet (pp_event_t pp_event)`

Returns the state on the associated event.

Parameters:

in	<i>pp_event</i>	Pointer to an associated event.
----	-----------------	---------------------------------

Return values:

<i>EV_RESET</i>	Event Reset.
<i>EV_SET</i>	Event Set.
<i>EV_INVALID</i>	Invalid Event.

`void R_OS_EventReset (pp_event_t pp_event)`

Clears the state on the associated event. Setting event to EV_RESET.

Parameters:

in	<i>pp_event</i>	Pointer to a associated event.
----	-----------------	--------------------------------

Returns:

none.

`void R_OS_EventSet (pp_event_t pp_event)`

Sets the state on the associated event outside of an interrupt service routine. Setting event to EV_SET.

Parameters:

in	<i>pp_event</i>	Pointer to an associated event.
----	-----------------	---------------------------------

Returns:

none.

`bool_t R_OS_EventSetFromIsr (pp_event_t pp_event)`

Sets the state on the associated event from inside an interrupt service routine. Setting event to EV_SET

Warning:

Function shall only be called from within an ISR routine

Parameters:

in	<i>pp_event</i>	Pointer to an associated event
----	-----------------	--------------------------------

Returns:

The function returns TRUE if the event object was successfully set. Otherwise, FALSE is returned

`bool_t R_OS_EventWait (pp_event_t pp_event, systime_t timeout)`

Blocks operation until one of the following occurs

A timeout occurs.

The associated event has been set.

Parameters:

in	<i>pp_event</i>	Pointer to an associated event.
in	<i>timeout</i>	Maximum time to wait for associated event to occur.

Returns:

The function returns TRUE if the event object was set, Otherwise, FALSE is returned

`void R_OS_Free (void ** pp_memory_to_free)`

Function to free allocated memory.

Parameters:

in	<i>p_memory_to_free</i>	Block of memory to free.
----	-------------------------	--------------------------

Returns:

None.

`uint32_t R_OS_GetTickCount (void)`

Gets ticks currently counted for task which calls it.

Warning:

Function can only be called when the scheduler is running

Returns:

The function returns the number of ticks counted.

`int32_t R_OS_GetVersion (st_os_abstraction_info_t * p_info)`

Obtains the version information from this module.

Parameters:

in	<i>p_info</i>	Structure containing module version information.
----	---------------	--------------------------------------------------

Returns:

The function returns 0

`void R_OS_KernelInit (void)`

Function to configure critical resources for the connected OS or scheduler, or configure an OS-Less sample.

Return values:

<i>NONE.</i>	
--------------	--

`void R_OS_Running (void)`

Function used to determine if the connected OS or scheduler has started.

Return values:

<i>TRUE</i>	Scheduler has started
<i>FALSE</i>	Scheduler has not started

`void R_OS_KernelStart (void)`

Function to enable the connected OS or scheduler, or configure an OS-Less sample.

Return values:

<i>NONE.</i>	
--------------	--

void R_OS_KernelStop (void)

Function to stop the connected OS or scheduler, or configure an OS-Less sample. Provided for completeness, may never be used. When powering down a system safely this function should be called.

Return values:

<i>NONE.</i>	
--------------	--

void* R_OS_Malloc (size_t *size*, e_memory_region_t *region*)

Allocates block of memory the length of "size".

Parameters:

in	<i>size</i>	Size of memory to allocate.
in	<i>region</i>	Region of memory to allocate from.

Returns:

Allocated memory

bool_t R_OS_MessageQueueClear (p_os_msg_queue_handle_t p_queue_handle)

Clear a message queue, resetting it to an empty state.

Parameters:

in	p_queue_handle	pointer to queue handle.
----	----------------	--------------------------

Returns:

The function returns TRUE if the event object was successfully cleared. Otherwise, FALSE is returned

bool_t R_OS_MessageQueueCreate (p_os_msg_queue_handle_t * pp_queue_handle, uint32_t queue_sz)

Create a Message Queue of length "queue_sz".

Parameters:

in	<i>queue_sz</i>	Maximum number of elements in queue.
in	<i>pp_queue_handle</i>	pointer to queue handle pointer.

Return values:

<i>TRUE</i>	The message queue was successfully created
<i>FALSE</i>	The message queue creation failed.

```
bool_t R_OS_MessageQueueDelete (p_os_msg_queue_handle_t * pp_queue_handle)
```

Delete a message queue. The message queue pointer argument will be set to NULL.

Parameters:

in	<i>pp_queue_handle</i>	pointer to queue handle pointer.
----	------------------------	----------------------------------

Returns:

The function returns TRUE if the event object was successfully deleted. Otherwise, FALSE is returned

```
bool_t R_OS_MessageQueueGet (p_os_msg_queue_handle_t p_queue, p_os_msg_t * pp_msg, uint32_t timeout, bool_t blocking)
```

Retrieve a message from a queue. Can only be called outside of an Interrupt Service Routine.

Parameters:

in	<i>p_queue</i>	pointer to queue handle.
out	<i>pp_msg</i>	pointer to message pointer. Pointer will point to NULL if no message and times out.
in	<i>timeout</i>	in system ticks.
in	<i>blocking</i>	true = block thread/task until message received. False = not blocking

Returns:

The function returns TRUE if the event object was successfully retrieved from the queue. Otherwise, FALSE is returned

```
bool_t R_OS_MessageQueuePut (p_os_msg_queue_handle_t p_queue_handle, p_os_msg_t p_message)
```

Put a message onto a queue. Can be called from both inside and outside of an Interrupt Service Routine.

Parameters:

in	<i>p_queue_handle</i>	pointer to queue handle.
in	<i>p_message</i>	pointer to message.

Returns:

The function returns TRUE if the event object was successfully added to the queue. Otherwise, FALSE is returned

```
void R_OS_MutexAcquire (p_mutex_t p_mutex)
```

Acquires possession of a Mutex, will context switch until free, with no timeout.

Parameters:

in	<i>p_mutex</i>	Mutex object to acquire.
----	----------------	--------------------------

Returns:

None.

```
void* R_OS_MutexCreate (void )
```

Creates a mutex and returns a pointer to it.

Return values:

<i>p_mutex</i>	Pointer to mutex created.
<i>NULL</i>	If mutex creation fails.

```
void R_OS_MutexDelete (pp_mutex_t pp_mutex)
```

Deletes a Mutex.

Parameters:

in	<i>pp_mutex</i>	Address of mutex object to delete, set to NULL when deleted.
----	-----------------	--------------------------------------------------------------

Returns:

None.

```
void R_OS_MutexRelease (p_mutex_t p_mutex)
```

Releases possession of a mutex.

Parameters:

in	<i>p_mutex</i>	Mutex object to release.
----	----------------	--------------------------

Returns:

None.

`bool_t R_OS_MutexWait (pp_mutex_t pp_mutex, uint32_t time_out)`

Attempts to claim mutex for 'timeout' length, will fail if not possible. If mutex passed is NULL, this function will create new mutex.

Parameters:

in	<i>pp_mutex</i>	Address of mutex object to acquire.
in	<i>time_out</i>	Length of Time to wait for mutex.

Return values:

<i>TRUE</i>	Mutex is acquired
<i>FALSE</i>	Wait Timed out, mutex not acquired.

`bool_t R_OS_SemaphoreCreate (p_semaphore_t p_semaphore, uint32_t count)`

Create a semaphore.

Parameters:

in	<i>p_semaphore</i>	Pointer to an associated semaphore.
in	<i>count</i>	The maximum count for the semaphore object. This value must be greater than zero

Return values:

<i>TRUE</i>	The semaphore object was successfully created.
<i>FALSE</i>	Semaphore not created.

`void R_OS_SemaphoreDelete (p_semaphore_t p_semaphore)`

Delete a semaphore, freeing any associated resources.

Parameters:

in	<i>p_semaphore</i>	Pointer to an associated semaphore.
----	--------------------	-------------------------------------

Returns:

None.

```
void R_OS_SemaphoreRelease (p_semaphore_t p_semaphore)
```

Release a semaphore, freeing it to be used by another task.

Parameters:

in	<i>p_semaphore</i>	Pointer to an associated semaphore.
----	--------------------	-------------------------------------

Returns:

None.

```
bool_t R_OS_SemaphoreWait (p_semaphore_t p_semaphore, systime_t timeout)
```

Blocks operation until one of the following occurs

A timeout occurs.

The associated semaphore has been set.

Parameters:

in	<i>p_semaphore</i>	Pointer to an associated semaphore.
in	<i>timeout</i>	Maximum time to wait for associated event to occur.

Return values:

<i>TRUE</i>	The semaphore object was successfully set.
<i>FALSE</i>	Semaphore not set.

```
int_t R_OS_SysLock (void )
```

Function to lock a critical section.

Warning:

This function must prevent the OS or scheduler from swapping context. This is often implemented by preventing system interrupts from occurring, and so pending any OS timer interruptions. Timing is critical, code protected by this function must be able to complete in the minimum time possible and never block.

Return values:

<i>1</i>	Critical Section entered
<i>0</i>	Object locked
<i>-1</i>	Error, neither action possible

void R_OS_SysReleaseAccess (void)

Function to release system mutex.

The OS Abstraction layer contains a system mutex. This function allows a user to release the mutex from system critical usage.

Returns:

None.

void R_OS_SysUnlock (void)

Function to unlock a critical section.

Warning:

This function releases the OS or scheduler to normal operation. Timing is critical, code proceeding this function must be able to complete in the minimum time possible and never block.

Returns:

None.

void R_OS_SysWaitAccess (void)

Function to acquire system mutex.

The OS Abstraction layer contains a system mutex. This function allows a user to obtain the mutex for system critical usage.

Returns:

None.

`os_task_t* R_OS_TaskCreate (const char_t * p_name, os_task_code_t task_code, void * p_params, size_t stack_size, int_t priority)`

Function to create a new task.

Parameters:

in	<i>p_name</i>	ASCII character representation for the name of the Task.
----	---------------	----------------------------------------------------------

Warning:

name string may be subject to length limitations. There is a security risk if the name is not bounded effectively in the implementation.

Parameters:

in	<i>task_code</i>	Function pointer to the implementation of the Task.
in	<i>p_params</i>	Structure to be used by the Task.
in	<i>stack_size</i>	Stack size for allocation to the Task.
in	<i>priority</i>	Task priority in system context.

Return values:

<i>os_task_t</i>	The task object.
------------------	------------------

`void R_OS_TaskDelete (os_task_t ** p_task)`

Function to delete a task.

Warning:

The target OS is responsible for verifying the Task is valid to delete.

Parameters:

in	<i>p_task</i>	the task object.
----	---------------	------------------

Return values:

<i>None.</i>	
--------------	--

`os_task_t* R_OS_TaskGetCurrentHandle (void)`

Gets current task.

Warning:

Function shall only be called when the scheduler is running

Parameters:

in	<i>none</i>	
----	-------------	--

Returns:

The function returns the current running task

`const char* R_OS_TaskGetCurrentName (void)`

Gets text name of current task.

Warning:

Function shall only be called when the scheduler is running

Parameters:

in	<i>none</i>	
----	-------------	--

Returns:

The function returns a pointer to the text name of the current task

int32_t R_OS_TaskGetPriority (uint32_t *task_id*)

Gets current task priority.

Warning:

Function shall only be called when the scheduler is running

Parameters:

in	<i>task_id</i>	desired Task
----	----------------	--------------

Returns:

The function returns the task priority of the specified uiTaskID
-1 if the uiTaskID can not be found

const char* R_OS_TaskGetState (const char * *p_task*)

Gets status information on selected task in human readable form.

Warning:

Function shall only be called when the scheduler is running

Parameters:

in	<i>p_task</i>	task name in human readable form.
----	---------------	-----------------------------------

Returns:

The function returns a character string that can be displayed on a console.

`bool_t R_OS_TaskResume (os_task_t * p_task)`

Function to cause a task to suspend and pass control back to the OS / scheduler.

Parameters:

in	<i>task</i>	the task object.
----	-------------	------------------

Return values:

<i>None.</i>	
--------------	--

`bool_t R_OS_TaskSetPriority (uint32_t task_id, uint32_t priority)`

Sets current task priority.

Warning:

Function shall only be called when the scheduler is running

Parameters:

in	<i>task_id</i>	desired task
in	<i>Priority</i>	desired priority

Returns:

true if priority is set
false if priority can not be set

`uint32_t R_OS_TasksGetNumber (void)`

Function to obtain total number of active tasks defined in the system, only attempted if the operating system is running.

Return values:

Number of tasks

`void R_OS_TaskSleep (uint32_t sleep_ms)`

Function to cause a task to suspend and pass control back to the OS / scheduler for a requested period.

Warning:

The time stated is a minimum, higher priority tasks may prevent this Task from being restored immediately.

Parameters:

in	<i>sleep_ms</i>	Time in ms (uint32 => max ~49 Days).
----	-----------------	--------------------------------------

Return values:

<i>None.</i>	
--------------	--

`void R_OS_TasksResumeAll (void)`

Resume all tasks, only attempted if the operating system is running.

Parameters:

None.	
-------	--

Return values:

None.	
-------	--

`void R_OS_TasksSuspendAll (void)`

Suspend all tasks, only attempted if the operating system is running.

Parameters:

None.	
-------	--

Return values:

None.	
-------	--

`bool_t R_OS_TaskSuspend (os_task_t * p_task)`

Function to cause a task to suspend and pass control back to the OS / scheduler.

Parameters:

in	<i>p_task</i>	the task object.
----	---------------	------------------

Return values:

None.	
-------	--

`void R_OS_TaskUsesFloatingPoint (void)`

Function to indicate to the OS that the current task uses floating point numbers.

`R_OS_TaskUsesFloatingPoint` should be called in the task that use FPU.

Return values:

NONE.	
-------	--

void R_OS_TaskYield (void)

Function to cause a task to suspend and pass control back to the OS / scheduler.

Return values:

<i>None.</i>	
--------------	--

5. Data Structure Documentation

5.1 st_os_abstraction_info_t Struct Reference

```
#include <r_os_abstraction_typedef.h>
```

Data Fields

```
union {  
    uint32_t full  
    struct {  
        uint16_t minor  
        uint16_t major  
    } sub  
} version  
uint32_t build  
const char * p_szdriver_name
```

Field Documentation

uint32_t build

Build Number Generated during the release

uint32_t full

Major + Minor combined as 1 uint32_t data member

uint16_t major

Version, modified by developer

uint16_t minor

Version, modified by Product Owner

const char* p_szdriver_name

struct { ... } sub

union { ... } version

The documentation for this struct was generated from the following file:

- **r_os_abstraction_typedef.h**

6. OS-Less OS Abstraction

The OS-less OS Abstraction is designed to provide some of the functionality of an OS to a non-OS environment. As it uses the common OS abstraction API, the task of porting between OS based and non-OS applications is simplified.

6.1 Supported Function API

The OS-less OS abstraction supports a reduced subset of the OS abstraction API. Table 6-1 below describes a list the OS abstraction functions and their status. Note that attempts to use unsupported functions will result in an “assert” handled error. These functions can be completed by the developer in an individual application specific way if portability is needed.

Function	Supported	Comments
R_OS_AbstractionLayerInit	✓	Implemented. Starts system timer if not already started.
R_OS_AbstractionLayerShutdown	✓	Implemented. Stops system timer if not already stopped.
R_OS_KernelInit	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_Running	✓	Implemented. Always returns TRUE.
R_OS_KernelStart	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_KernelStop	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_InitMemManager	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskCreate	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskDelete	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskSleep	✓	Wait for specified number of OS timer ticks.
R_OS_TaskYield	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskSuspend	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskResume	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TasksSuspendAll	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TasksResumeAll	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TasksGetNumber	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskUsesFloatingPoint	✗	Returns without doing anything. Does not call assert function
R_OS_TaskGetPriority	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskSetPriority	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskGetCurrentHandle	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskGetCurrentName	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_TaskGetState	✗	Not implemented. Calls assert function R_OS_AssertCalled
R_OS_SysLock	✓	Disables Interrupts
R_OS_SysUnlock	✓	Enables Interrupts
R_OS_SysWaitAccess	✗	Returns without doing anything. Does not call assert function
R_OS_SysReleaseAccess	✗	Returns without doing anything. Does not call assert function
R_OS_GetTickCount	✓	Returns current system tick count.
R_OS_AssertCalled	✓	Places execution into an infinite loop after recording file and line number. Can be used for debug purposes.
R_OS_Malloc	✓	Ignores region parameter and uses system malloc
R_OS_Free	✓	Uses system free

Function	Supported	Comments
R_OS_SemaphoreCreate	✓	
R_OS_SemaphoreDelete	✓	
R_OS_SemaphoreWait	✓	
R_OS_SemaphoreRelease	✓	
R_OS_MutexCreate	✓	The #define R_OSFREE_MAX_MUTEXES defines the number of mutexes available in the system. This can be adjusted to suit the application.
R_OS_MutexDelete	✓	
R_OS_MutexAcquire	✓	
R_OS_MutexRelease	✓	
R_OS_MutexWait	✓	
R_OS_EnterCritical	✓	Disables Interrupts
R_OS_ExitCritical	✓	Enables Interrupts
R_OS_MessageQueueCreate	✓	
R_OS_MessageQueuePut	✓	
R_OS_MessageQueueGet	✓	
R_OS_MessageQueueClear	✓	
R_OS_MessageQueueDelete	✓	
R_OS_EventCreate	✓	The #define R_OSFREE_MAX_EVENTS defines the number of events available in the system. This can be adjusted to suit the application.
R_OS_EventDelete	✓	
R_OS_EventSet	✓	
R_OS_EventReset	✓	
R_OS_EventGet	✓	
R_OS_EventWait	✓	
R_OS_EventSetFromIsr	✓	
R_OS_GetVersion	✓	

Table 6-1 : API functions in OS-less Abstraction

6.2 Connections to external non-API components

6.2.1 Compiler Abstraction

The OS abstraction layer uses the compiler abstraction in order to access simple assembly commands, as defined in the API in "r_compiler_abstraction_api.h".

6.2.2 System Timer

The OS abstraction middleware uses the OSTM timer peripheral to create the system tick functionality. An Interrupt Service Routine (ISR) function, os_abstraction_isr, is called when the timer counter overflows every millisecond, and this increments the system tick counter.

This is achieved by including the ostm driver, using Smart Configurator to set the OSTM peripheral to the correct channel, interval and ISR function.

7. FreeRTOS OS Abstraction

7.1 Supported Function API

The FreeRTOS OS Abstraction is designed to simplify the task of porting application code between Operating Systems.

The FreeRTOS OS abstraction implements the OS abstraction API as a layer above the FreeRTOS instance in the application project. Table 7-1 below describes a list of the OS abstraction functions and their status.

Function	Supported FreeRTOS: Amazon	Comments
R_OS_AbstractionLayerInit	✓	Calls R_OS_KernelInit
R_OS_AbstractionLayerShutdown	✓	Calls R_OS_KernelStop
R_OS_KernelInit	✓	Calls R_OS_InitMemManager, creates main_task then calls R_OS_KernelStart
R_OS_Running	✓	Implemented. Implemented state true(yes), false(no)
R_OS_KernelStart	✓	
R_OS_KernelStop	✓	
R_OS_InitMemManager	✓	Initialise heap in freeRTOS
R_OS_TaskCreate	✓	
R_OS_TaskDelete	✓	
R_OS_TaskSleep	✓	
R_OS_TaskYield	✓	
R_OS_TaskSuspend	✓	
R_OS_TaskResume	✓	
R_OS_TasksSuspendAll	✓	
R_OS_TasksResumeAll	✓	
R_OS_TasksGetNumber	✓	
R_OS_TaskUsesFloatingPoint	✓	
R_OS_TaskGetPriority	✓	
R_OS_TaskSetPriority	✓	
R_OS_TaskGetCurrentHandle	✓	
R_OS_TaskGetCurrentName	✓	
R_OS_TaskGetState	✓	
R_OS_SysLock	✓	
R_OS_SysUnlock	✓	
R_OS_SysWaitAccess	✓	
R_OS_SysReleaseAccess	✓	
R_OS_GetTickCount	✓	Returns current system tick count.
R_OS_AssertCalled	✓	Places execution into an infinite loop after recording file and line number data to console. Can be used for debug purposes.

Function	Supported FreeRTOS or Amazon	Comments
R_OS_Malloc	✓	FreeRTOS variant Supports memory region selection allowing application to select preferred memory region. Amazon Variant Does not support memory region selection.
R_OS_Free	✓	
R_OS_SemaphoreCreate	✓	
R_OS_SemaphoreDelete	✓	
R_OS_SemaphoreWait	✓	
R_OS_SemaphoreRelease	✓	
R_OS_MutexCreate	✓	
R_OS_MutexDelete	✓	
R_OS_MutexAcquire	✓	
R_OS_MutexRelease	✓	
R_OS_MutexWait	✓	
R_OS_EnterCritical	✓	
R_OS_ExitCritical	✓	
R_OS_MessageQueueCreate	✓	
R_OS_MessageQueuePut	✓	
R_OS_MessageQueueGet	✓	
R_OS_MessageQueueClear	✓	
R_OS_MessageQueueDelete	✓	
R_OS_EventCreate	✓	
R_OS_EventDelete	✓	
R_OS_EventSet	✓	
R_OS_EventReset	✓	
R_OS_EventGet	✓	
R_OS_EventWait	✓	
R_OS_EventSetFromIsr	✓	
R_OS_GetVersion	✓	

Table 7-1 : API functions in FreeRTOS OS Abstraction

7.2 Task Priorities

The FreeRTOS OS abstraction has a header file “r_task_priority.h” which is used to define the priorities of system tasks, such as the main task, console, idle task etc.

7.3 Connections to external non-API components

7.3.1 Compiler Abstraction

The OS abstraction layer uses the compiler abstraction in order to access simple assembly commands, as defined in the API in "r_compiler_abstraction_api.h".

7.3.2 System Timer

The OS abstraction middleware uses the OSTM timer peripheral to create the system tick functionality. An Interrupt Service Routine (ISR) function, `os_abstraction_isr`, is called when the timer counter overflows every millisecond.

This is achieved by including the ostm driver, using Smart Configurator to set the OSTM peripheral to the correct channel, interval and ISR function.

7.3.3 FreeRTOS

The OS abstraction layer uses freeRTOS to implement the functionality. As such it includes the following headers, which should be made available in the project

```
#include "FreeRTOS.h"
#include "FreeRTOSconfig.h"
#include "semphr.h"
#include "queue.h"
#include "task.h"
```

7.3.4 Configuring Memory Regions

Memory can be defined allowing support of multiple non adjacent (non-contiguous) memory regions.

FreeRTOS OS abstraction layer includes an enhanced version of the default heap5.c memory module to support the selection of which region is used in the `R_OS_Malloc` function. When using the FreeRTOS module, the `e_memory_region_t` parameter is used to specify which region is preferred for the allocation.

Certified Amazon FreeRTOS OS supports the specification of memory regions, but does not support the selection of which region `R_OS_Malloc()` uses. Amazon certification prohibits the modification of core software. When using the certified Amazon module, the `e_memory_region_t` parameter is ignored.

To configure the memory regions (used in both variants) create the `e_memory_region_t` table in the following file: `generate/system/inc/r_typedefs.h` and define the regions in any .c file (ie `main.c`).

Example `r_typedefs.h`

```
#define REPLACE_MEMORY_REGION_ENUM (1) /* Override the default implementation */
typedef enum
{
    R_MEMORY_REGION_DEFAULT = 0,
    R_MEMORY_REGION_NEW,
} e_memory_region_t;
```

Example `main.c`

```
#include "FreeRTOS.h"
#include "r_compiler_abstraction_api.h"
#include "r_os_abstraction_api.h"

HeapRegion_t xHeapRegions[] =
{
    {( uint8_t * ) 0x80080000UL, 0x00060000UL}, /* R_MEMORY_REGION_DEFAULT */
    {( uint8_t * ) 0x800E0000UL, 0x00020000UL}, /* R_MEMORY_REGION_NEW */
    {( uint8_t * ) 0x00000000UL, 0x00000000UL}, /* Terminates the array */
};
```

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Revision History

Rev.	Date	Description	
		Page	Summary
3.10	21/03/2019	All	Created document to align with OS Abstraction layer V.3.10
3.20	25/07/2019	All	Filename change
3.50	25/06/2020	4.2	Added Known limitation, align document version to API revision

General Precautions in the Handling of Microprocessing Unit and Microcontroller Unit Products

The following usage notes are applicable to all Microprocessing unit and Microcontroller unit products from Renesas. For detailed usage notes on the products covered by this document, refer to the relevant sections of the document as well as any technical updates that have been issued for the products.

1. Precaution against Electrostatic Discharge (ESD)

A strong electrical field, when exposed to a CMOS device, can cause destruction of the gate oxide and ultimately degrade the device operation. Steps must be taken to stop the generation of static electricity as much as possible, and quickly dissipate it when it occurs. Environmental control must be adequate. When it is dry, a humidifier should be used. This is recommended to avoid using insulators that can easily build up static electricity.

Semiconductor devices must be stored and transported in an anti-static container, static shielding bag or conductive material. All test and measurement tools including work benches and floors must be grounded. The operator must also be grounded using a wrist strap. Semiconductor devices must not be touched with bare hands. Similar precautions must be taken for printed circuit boards with mounted semiconductor devices.

2. Processing at power-on

The state of the product is undefined at the time when power is supplied. The states of internal circuits in the LSI are indeterminate and the states of register settings and pins are undefined at the time when power is supplied. In a finished product where the reset signal is applied to the external reset pin, the states of pins are not guaranteed from the time when power is supplied until the reset process is completed. In a similar way, the states of pins in a product that is reset by an on-chip power-on reset function are not guaranteed from the time when power is supplied until the power reaches the level at which resetting is specified.

3. Input of signal during power-off state

Do not input signals or an I/O pull-up power supply while the device is powered off. The current injection that results from input of such a signal or I/O pull-up power supply may cause malfunction and the abnormal current that passes in the device at this time may cause degradation of internal elements. Follow the guideline for input signal during power-off state as described in your product documentation.

4. Handling of unused pins

Handle unused pins in accordance with the directions given under handling of unused pins in the manual. The input pins of CMOS products are generally in the high-impedance state. In operation with an unused pin in the open-circuit state, extra electromagnetic noise is induced in the vicinity of the LSI, an associated shoot-through current flows internally, and malfunctions occur due to the false recognition of the pin state as an input signal become possible.

5. Clock signals

After applying a reset, only release the reset line after the operating clock signal becomes stable. When switching the clock signal during program execution, wait until the target clock signal is stabilized. When the clock signal is generated with an external resonator or from an external oscillator during a reset, ensure that the reset line is only released after full stabilization of the clock signal. Additionally, when switching to a clock signal produced with an external resonator or by an external oscillator while program execution is in progress, wait until the target clock signal is stable.

6. Voltage application waveform at input pin

Waveform distortion due to input noise or a reflected wave may cause malfunction. If the input of the CMOS device stays in the area between V_{IL} (Max.) and V_{IH} (Min.) due to noise, for example, the device may malfunction. Take care to prevent chattering noise from entering the device when the input level is fixed, and also in the transition period when the input level passes through the area between V_{IL} (Max.) and V_{IH} (Min.).

7. Prohibition of access to reserved addresses

Access to reserved addresses is prohibited. The reserved addresses are provided for possible future expansion of functions. Do not access these addresses as the correct operation of the LSI is not guaranteed.

8. Differences between products

Before changing from one product to another, for example to a product with a different part number, confirm that the change will not lead to problems. The characteristics of a microprocessing unit or microcontroller unit products in the same group but having a different part number might differ in terms of internal memory capacity, layout pattern, and other factors, which can affect the ranges of electrical characteristics, such as characteristic values, operating margins, immunity to noise, and amount of radiated noise. When changing to a product with a different part number, implement a system-evaluation test for the given product.

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